

Enforcing the GDPR

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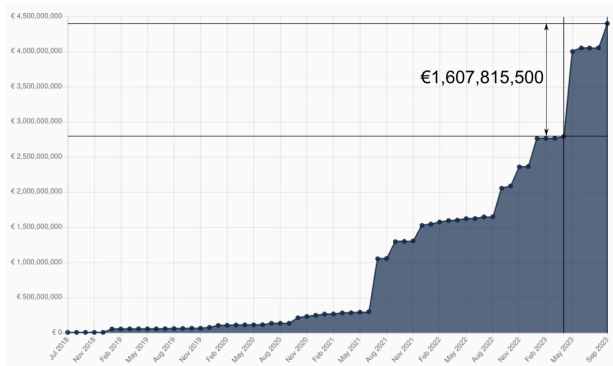
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More privacy thanks to the GDPR?

Since the time our paper was written, over **€1.5 billion** in fines have been imposed on entities violating the GDPR.



Source: www.enforcementtracker.com

How can we do better?

Be proactive!

How can we do better?

Be proactive!



Privacy by Design

The 7 Foundational Principles

Implementation and Mapping of Fair Information Practices

Cavoukian, A. (2009). *Privacy by Design*.

How can we do better?

Be proactive!



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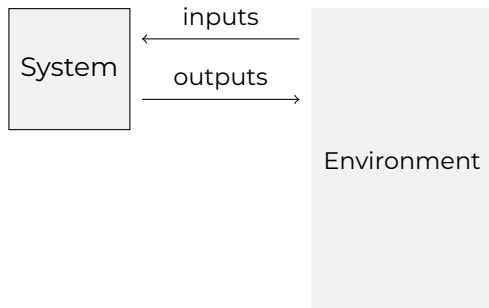
Cavoukian, A. (2009). *Privacy by Design*.

Data protection by design and by default

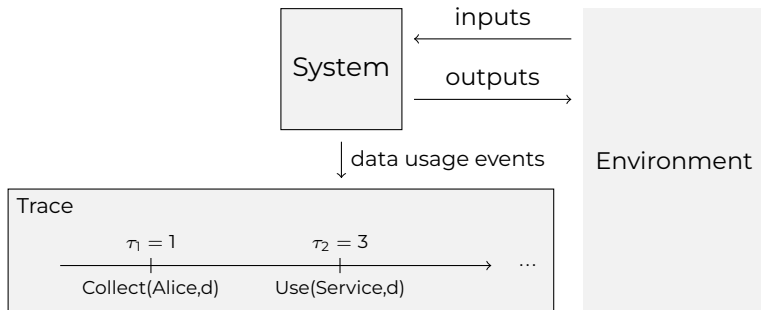
*Taking into account the state of the art, the cost of implementation and the nature, scope, context and purposes of processing [...] **the controller shall [...] implement appropriate technical and organisational measures [...] in order to meet the requirements of this Regulation** and protect the rights of data subjects.*

— GDPR, art. 25(1)

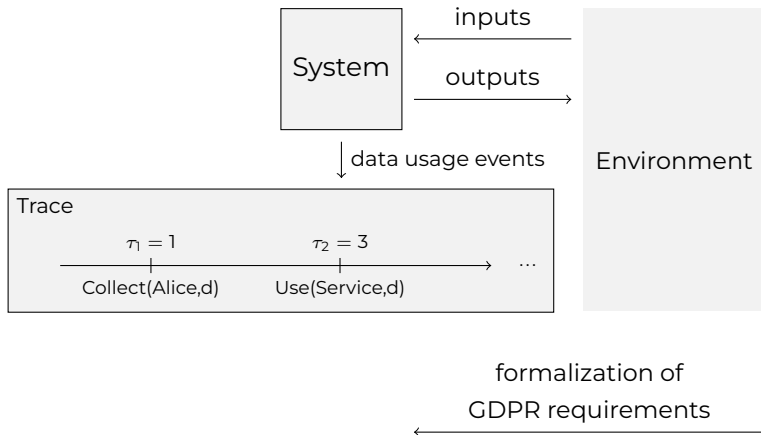
This paper: PbD using Runtime Enforcement



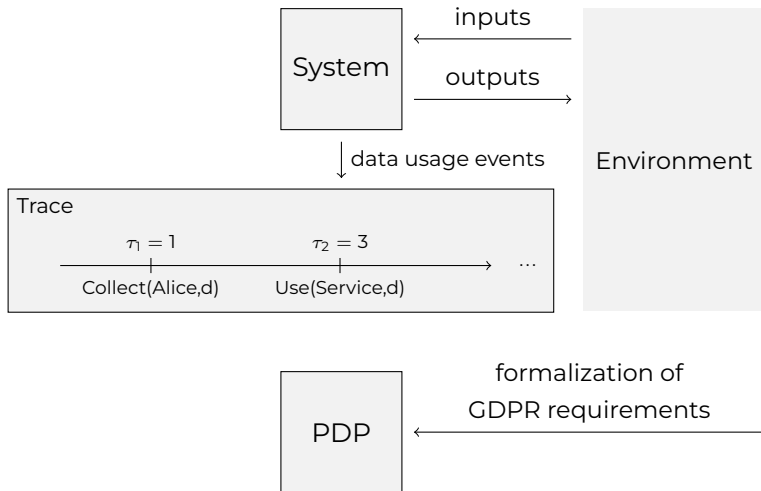
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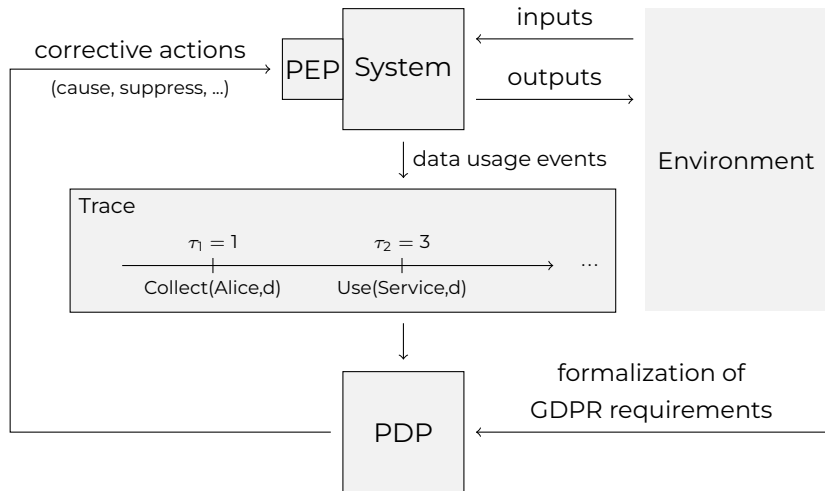
This paper: PbD using Runtime Enforcement



This paper: PbD using Runtime Enforcement



This paper: PbD using Runtime Enforcement



This paper: Runtime Enforcement of the GDPR

Challenges



Contributions



This paper: Runtime Enforcement of the GDPR

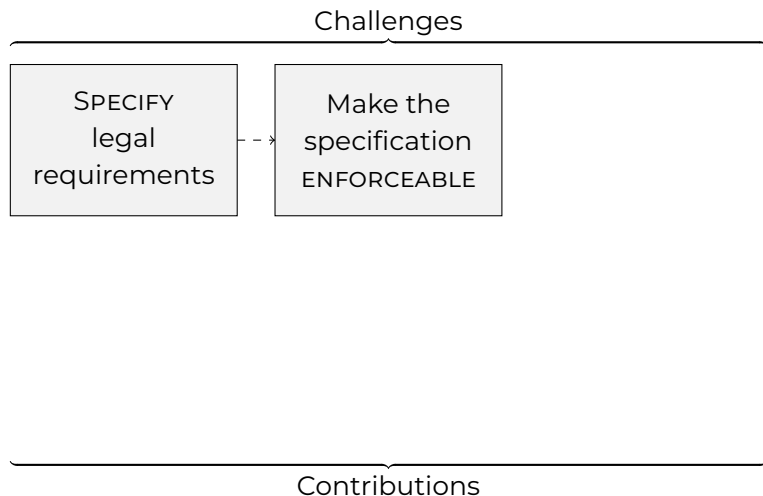
Challenges

The diagram features a light gray rectangular box on the left containing the text 'SPECIFY legal requirements'. Above this box is a horizontal line with a central upward-pointing curly bracket, labeled 'Challenges'. Below the entire diagram area is another horizontal line with a central downward-pointing curly bracket, labeled 'Contributions'.

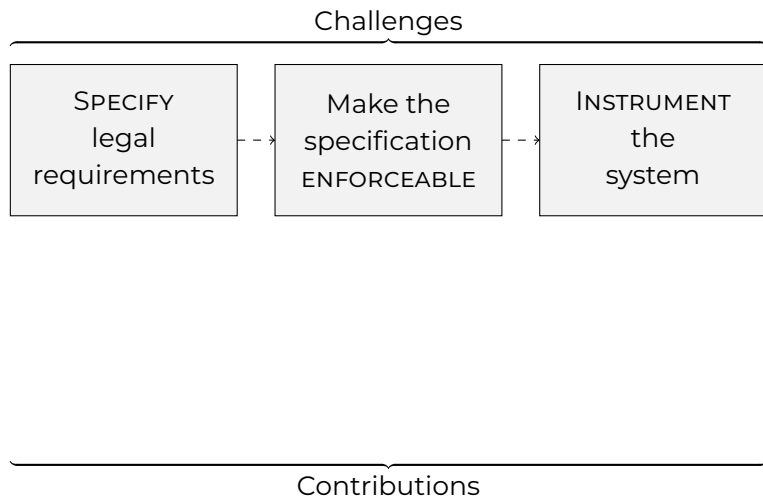
SPECIFY
legal
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Contributions

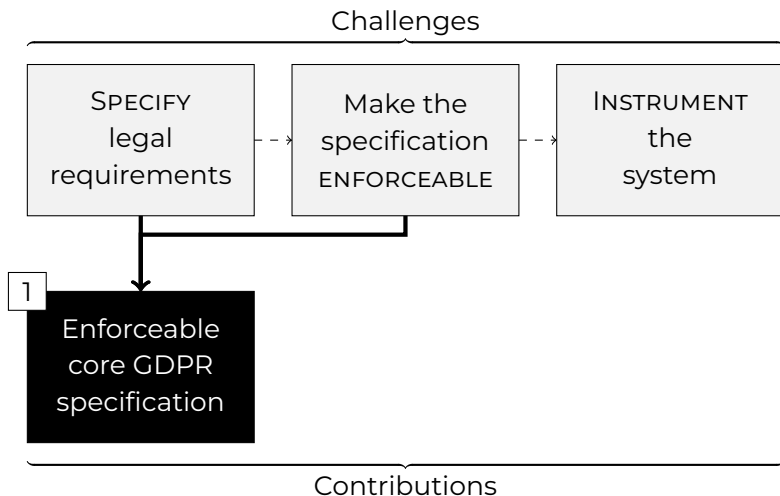
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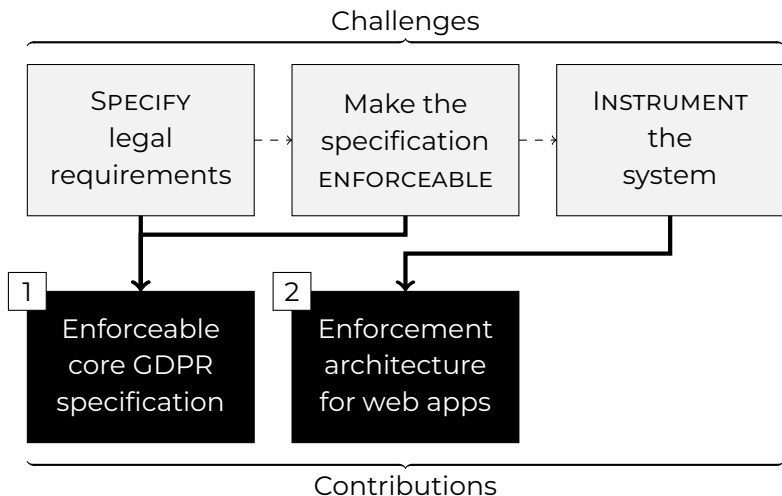
This paper: Runtime Enforcement of the GDPR



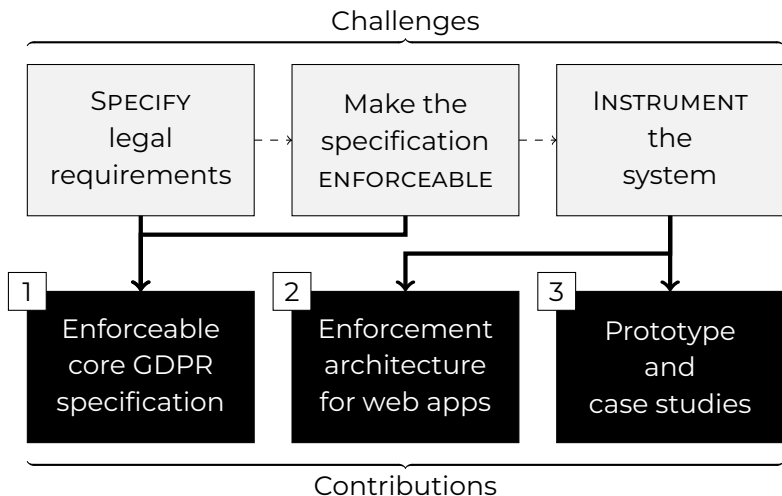
This paper: Runtime Enforcement of the GDPR



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This paper: Runtime Enforcement of the GDPR



1

Enforceable
core GDPR
specification

Previous work: Monitorable GDPR specification



Article	MFOTL Formula
5(1)(c) Data minimisation	collect(<i>data</i> , <i>dataid</i> , <i>dsid</i>) IMPLIES EVENTUALLY use(<i>data</i> , <i>dataid</i> , <i>dsid</i>)
5(1)(e) Storage limitation	collect(<i>data</i> , <i>dataid</i> , <i>dsid</i>) IMPLIES EVENTUALLY delete(<i>data</i> , <i>dataid</i> , <i>dsid</i>)
6(1) Lawful processing	use(<i>data</i> , <i>dataid</i> , <i>dsid</i>) IMPLIES ONCE (ds.consents(<i>dsid</i> , <i>data</i>) OR legal_grounds(<i>dsid</i> , <i>data</i>))
7(3) Consent	use(<i>data</i> , <i>dataid</i> , <i>dsid</i>) IMPLIES (ONCE legal_grounds(<i>dsid</i> , <i>data</i>) OR (NOT ds.revoke(<i>dsid</i> , <i>data</i>) SINCE ds.consents(<i>dsid</i> , <i>data</i>)))

Arfelt, E. et al. (2019). Monitoring the GDPR. *ESORICS'19*.

- Focuses on a simple specification of **data subject rights**
- Uses **Metric First-Order Temporal Logic** (MFOTL)
- **Monitors** traces for GDPR violations

Previous work: MFOTL enforcement



Algorithm 1. Function enf

```
function enf( $\varphi, \sigma, v$ )
  if  $\varphi = r(t_1, \dots, t_n), r \in \text{Sup}$  then
    return  $(\{(r, (v(t_1), \dots, v(t_n)))\}, \emptyset)$ 
  else if  $\varphi = \neg r(t_1, \dots, t_n), r \in \text{Cau}$  then
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- Supports **suppressable** and **causable** events
- Characterizes an **enforceable fragment** of MFOTL

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- Supports **suppressable** and **causable** events
- Characterizes an **enforceable fragment** of MFOTL
- Provides a PDP (**EnfPoly** tool)
- Two main inputs: **signature** + **MFOTL formula**

GDPR signature for enforcement

Event	Description	Type
	The data subject ds :	
$DSConsent(ds, prp, d)$	gives consent to use data d for purpose prp	
$DSRevoke(ds, prp, d)$	revokes consent given to use data d for purpose prp	
...
$DSErase(ds, d)$	requests erasure of data d	
...
	The application:	
$Collect(ds, d)$	collects data d of data subject ds	
$Use(prp, d)$	uses data d for purpose prp	Ⓢ
...
$Erase(d)$	erases data d	ⓐ
...
$LegalGround(grd, d)$	application claims legal ground grd for using data d	

ⓐ = Causable; Ⓢ = Suppressable; unmarked events are only-observable

Formalization of core GDPR requirements

We cover 6 core requirements / data subject rights:

- Purpose-based usage
- Right to access
- Right to rectification
- Right to erasure
- Right to restriction
- Right to object

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Core GDPR requirements using MFOTL

Example: **Purpose-based usage**

$$\alpha \hat{S} \beta := \alpha S (\alpha \wedge \beta)$$

$$\begin{aligned} \varphi_{\text{Purp}} = & \exists prp, d, ds. \text{Use}(prp, d) \wedge \blacklozenge \text{Collect}(ds, d) \\ & \wedge \neg((\neg \text{DSRevoke}(ds, prp, d) \hat{S} \text{DSConsent}(ds, prp, d)) \\ & \vee (\exists grd. \blacklozenge \text{LegalGround}(grd, d))) \end{aligned}$$

Purpose-based usage is violated when

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Purpose-based usage is violated when **some data d is used for purpose prp**

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Purpose-based usage is violated when some data d is used for purpose prp , **data d was collected from data subject ds**

Core GDPR requirements using MFOTL

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Purpose-based usage is violated when some data d is used for purpose prp , data d **was** collected from data subject ds **and neither**

Core GDPR requirements using MFOTL

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Core GDPR requirements using MFOTL (II)

Example: **Right to erasure**

$$\varphi_{\text{Erasure}} = \exists ds, ut. \text{DSErase}(ds, d) \wedge \blacklozenge \text{Collect}(ds, d) \wedge \neg \text{Erase}(d)$$

The right to erasure is violated when

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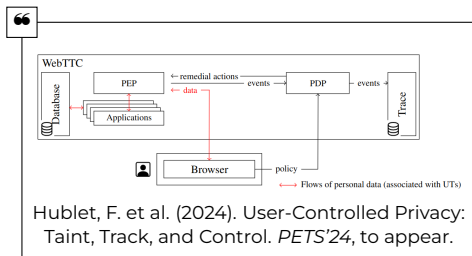
Enforceability

- Our formalization is **enforceable**, being in the enforceable fragment from Hublet et al. (2022)
- Enforcement through suppression / causation
 - Purpose based-usage: suppression of Use
 - Right to erasure: causation of $Erase$

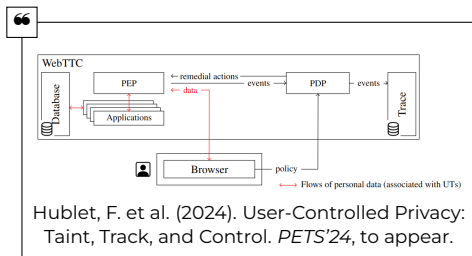
2

Enforcement
architecture
for web apps

Previous work: WebTTC



Previous work: WebTTC

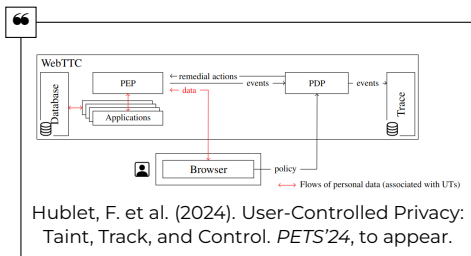


- Enforces **MFOTL privacy policies** in web applications by observing the following events:

Event	Description	Type
$In(u, d)$	Data subject ds inputs a value with UT d	
$Out(u, prp, d)$	Value with UT d is output to data subject ds for purpose prp	Ⓢ
$Itf(d, o)$	Input with UT d interferes with output o	

Ⓢ = Suppressable; unmarked events are only-observable

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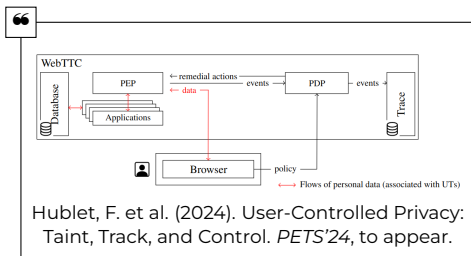
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$Out(u, prp, d)$	Value with UT d is output to data subject ds for purpose prp	Ⓢ
$Intf(d, o)$	Input with UT d interferes with output o	

Ⓢ = Suppressable; unmarked events are only-observable

- Python-like language** with dynamic Information Flow Control

Previous work: WebTTC



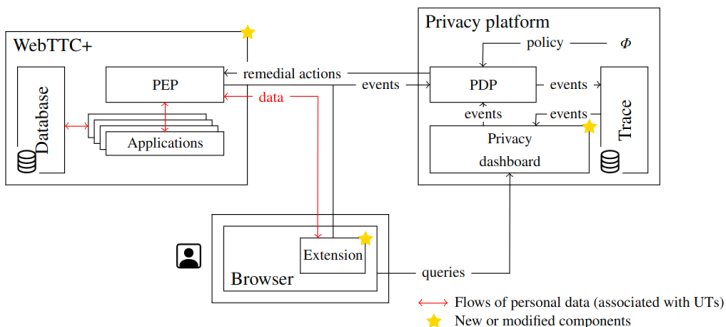
- Enforces **MFOTL privacy policies** in web applications by observing the following events:

Event	Description	Type
$In(u, d)$	Data subject ds inputs a value with UT d	
$Out(u, prp, d)$	Value with UT d is output to data subject ds for purpose prp	Ⓢ
$Intf(d, o)$	Input with UT d interferes with output o	

Ⓢ = Suppressable; unmarked events are only-observable

- Python-like language** with dynamic Information Flow Control
- Associates **unique taints** (UTs) to inputs and propagates them

Our enforcement architecture: WebTTC+

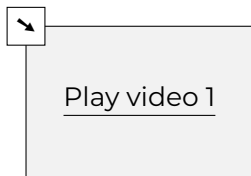


Instrumentation

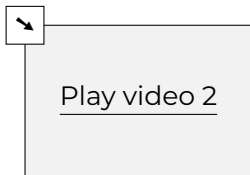
- Events are emitted at input or output time, on user request (DS...) or on PDP request (causable events)

Event	Type	When emitted?	By...
$DSConsent(ds, prp, d)$		Input time / on user request	Extension / Dashboard
$DSRevoke(ds, prp, d)$		On user request	Dashboard
$DSErase(ds, d)$		On user request	Dashboard
$Collect(ds, d, sp)$		Input time	Extension
$Use(prp, d)$	Ⓢ	Output time	PEP
$Erase(d)$	Ⓞ	On enforcer request	PEP

Instrumentation: input time



Instrumentation: user requests



3

Prototype
and
case studies

Prototype

Includes:

- WebTTC+
- Privacy platform including EnfPoly wrapper
- PoC browser extension
- EnfPoly (Hublet et al. 2022)

~ 3k lines of code on top of WebTTC and EnfPoly

Evaluation

We ported three applications from previous work (Wang et al. 2019, Yang et al. 2016).

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RQ2: How much runtime overhead does WebTTC+ incur
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Research questions:

- RQ1: Can realistic web applications be developed in WebTTC+?
How is the size of the code base impacted?
- RQ2: How much runtime overhead does WebTTC+ incur compared to a baseline without enforcement?
- RQ3: What share of the GDPRs provisions does our implementation effectively enforce?

Evaluation results

RQ1: Functionality preserved with low code overhead $\sim 10\%$

Evaluation results

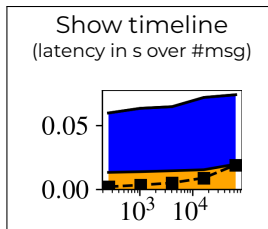
RQ1: Functionality preserved with low code overhead $\sim 10\%$

RQ2: Runtime overhead within 1 order of magnitude of baseline, latency < 75 ms, usability is preserved

Evaluation results

RQ1: Functionality preserved with low code overhead $\sim 10\%$

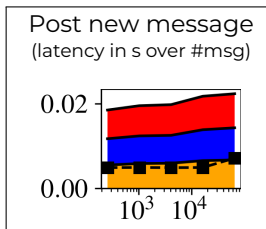
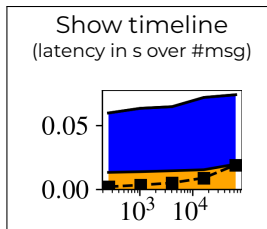
RQ2: Runtime overhead within 1 order of magnitude of baseline, latency < 75 ms, usability is preserved



Evaluation results

RQ1: Functionality preserved with low code overhead $\sim 10\%$

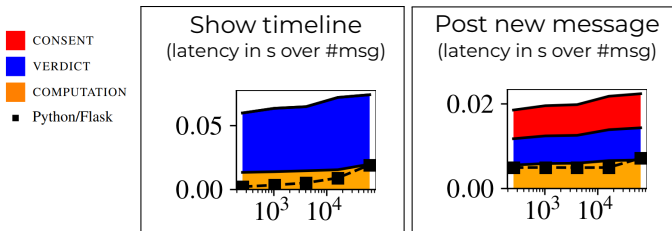
RQ2: Runtime overhead within 1 order of magnitude of baseline, latency < 75 ms, usability is preserved



Evaluation results

RQ1: Functionality preserved with low code overhead $\sim 10\%$

RQ2: Runtime overhead within 1 order of magnitude of baseline, latency < 75 ms, usability is preserved



RQ3: About $2/3$ of the fines imposed until May 2023 are related to at least one of the articles addressed by our approach

Conclusion and future work

- First work to provide an enforceable specification + enforcement architecture for a core of the GDPR

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Contributions



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Contributions

Enforceable
core GDPR
specification

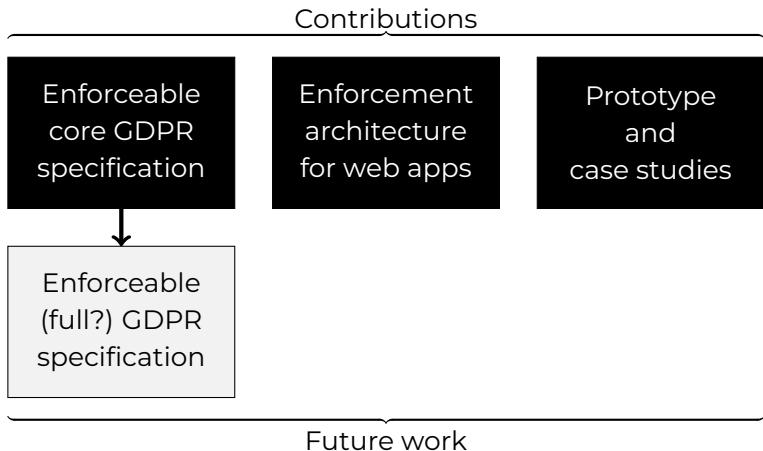
Enforcement
architecture
for web apps

Prototype
and
case studies

Future work

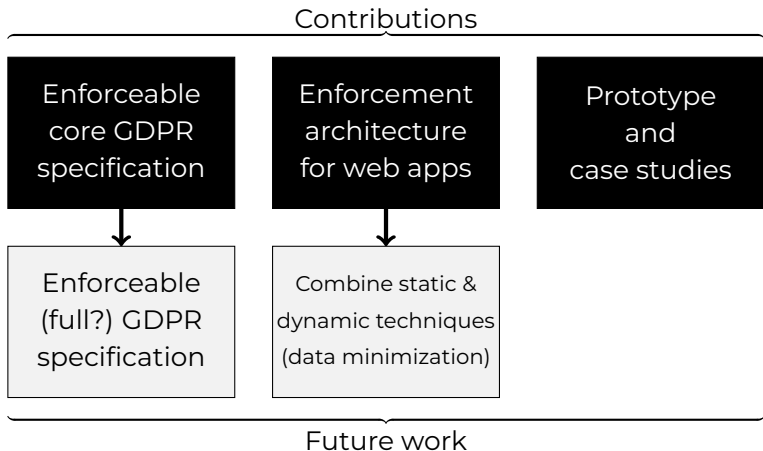
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